

# PNS FOR CAN

# PLUG-AND-SECURE COMMUNICATION FOR CAN

# Plug-and-Secure Communication for CAN Motivation

## Facts

- ▶ Current trends (e.g. Cloud/Internet connectivity) lead to novel & serious security threats
- ▶ Today's CAN networks are often hardly secured
- ▶ Cryptographic methods may help (e.g. message auth.)

## However

- ▶ Key agreement and distribution is **not** a solved or trivial problem  
Reasons: security, effort, computational complexity, price
- ▶ Keys have not been attacked – simply because they **did not exist**



## Our Idea: Plug-and-Secure

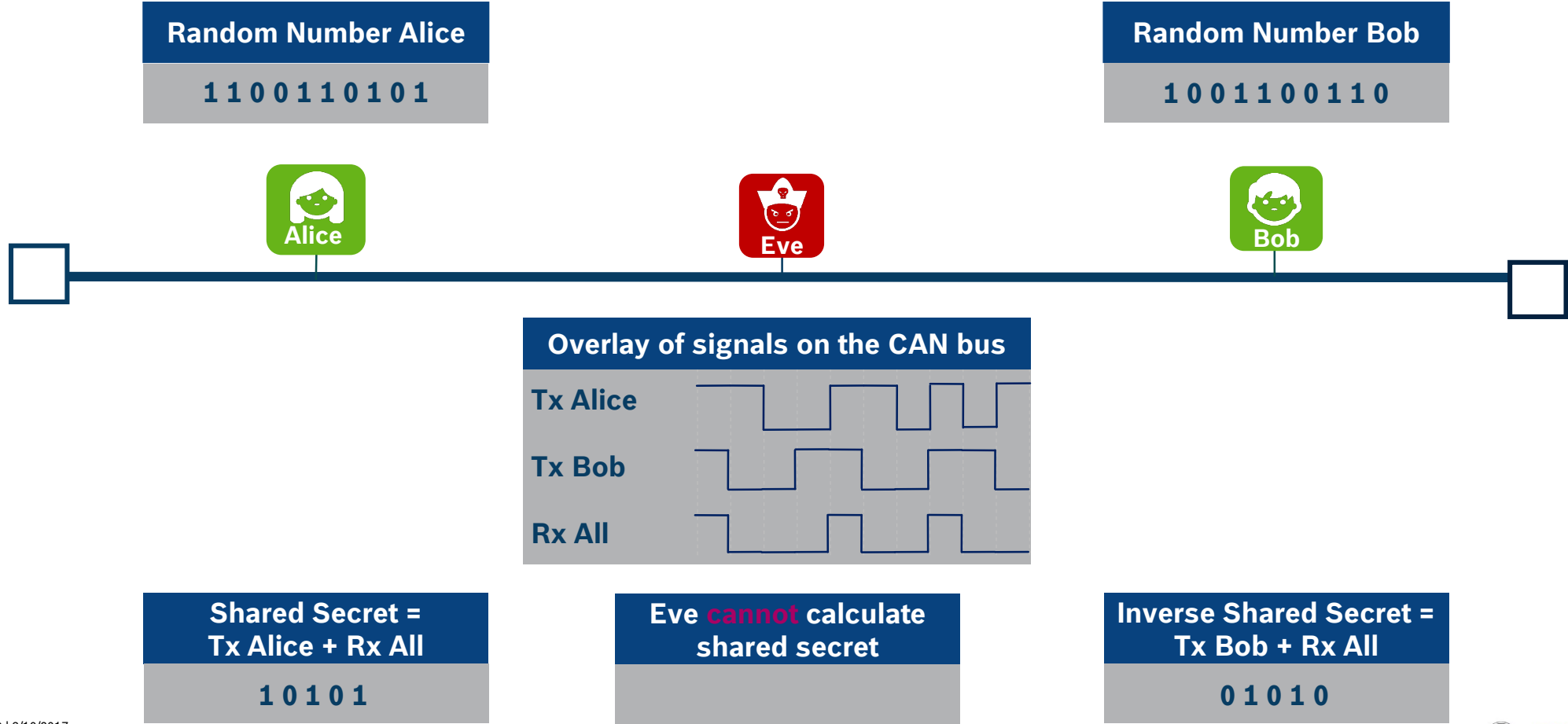
A novel approach for completely automated & secure key establishment of very low complexity for CAN networks (“plug-and-secure”)

Especially suitable against software-based & remote attack scenarios

Basic Idea: Exploit special properties of CAN bus (dominant / recessive bits)

# Plug-and-Secure Communication for CAN

## Basic idea



# Plug-and-Secure Communication for CAN Details

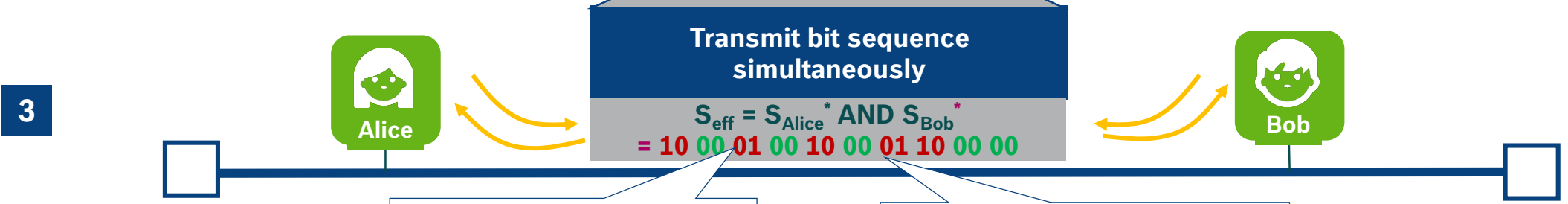
Alice	Bob	Bus
0	0	0
0	1	0
1	0	0
1	1	1

**1** Generate a random bit string of length N  
 $S_{Alice} = 1100110101$

Generate a random bit string of length N  
 $S_{Bob} = 1001100110$

**2** Replace: 0 → 01  
 1 → 10  
 $S_{Alice}^* = 1010010110100110$

Replace: 0 → 01  
 1 → 10  
 $S_{Bob}^* = 100101101001011001$



'10' or '01' = both users have transmitted identical bits

'00' = both users have transmitted different bits

**4**  $S_{Alice} = \cancel{X}1\cancel{X}0\cancel{X}1\cancel{X}0\cancel{X}1$   
 10101

Discard bits corresponding to '01 or '10' in  $S_{eff}$  in initial bit sequences  $S_{Alice} / S_{Bob}$

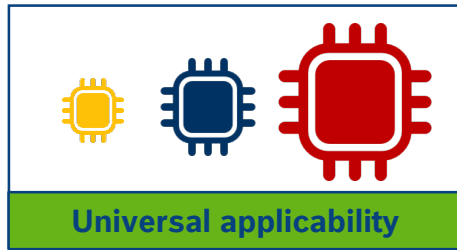
$S_{Bob} = \cancel{X}0\cancel{X}1\cancel{X}0\cancel{X}1\cancel{X}0$   
 01010

Inverse sequences = shared secret

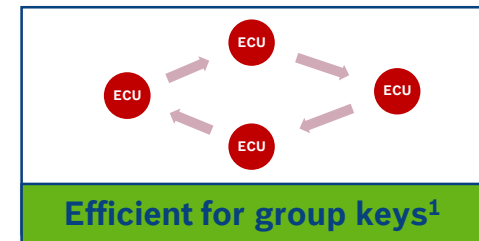


# Plug-and-Secure Communication for CAN

## Major benefits

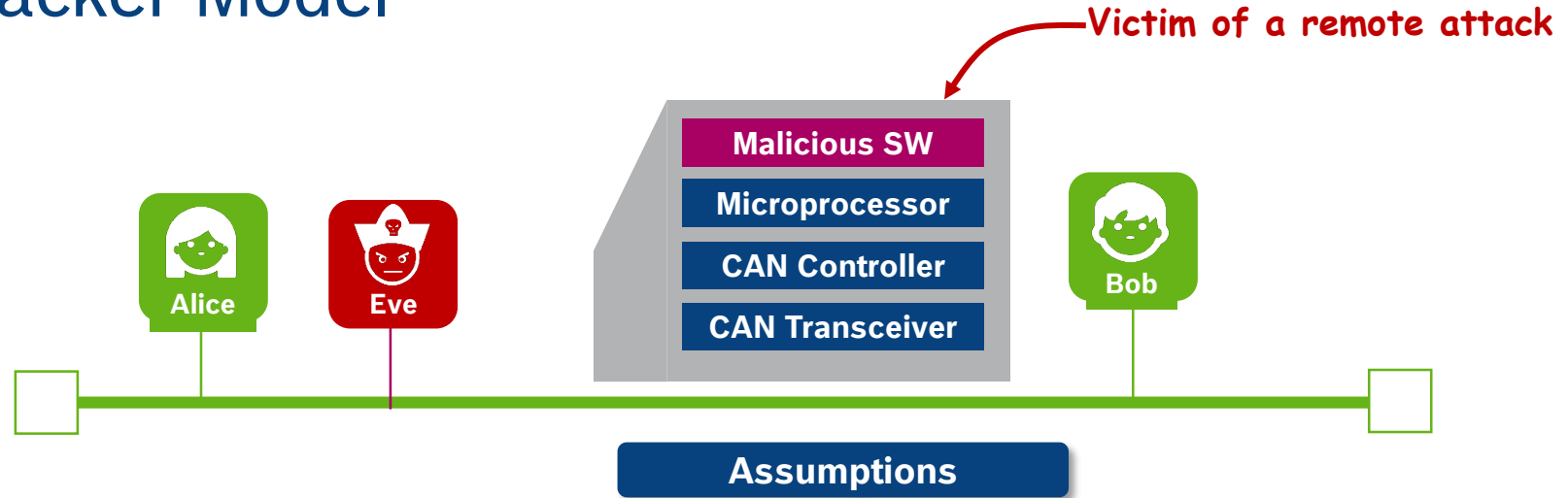


## Plug-and-Secure Communication for CAN



➔ Seamless integration into CAN ecosystem

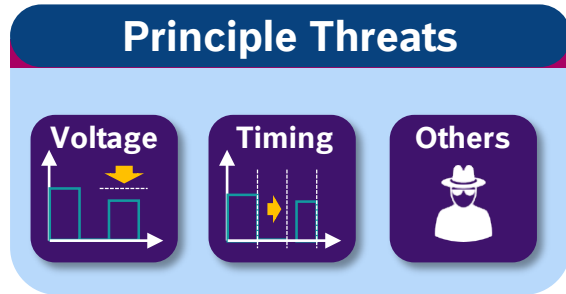
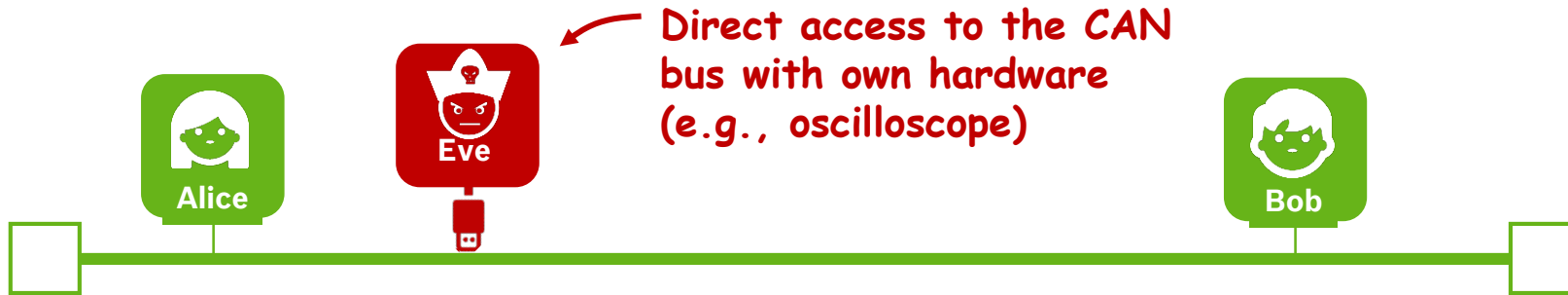
# Plug-and-Secure Communication for CAN Remote Attacker Model



- 1** Eve is using standard HW with modified (malicious) SW
- 2** Eve may eavesdrop on all messages exchanged on the CAN bus
- 3** Eve may inject arbitrary bits on the CAN bus (via the CAN transceiver)

**Highly relevant attacker model due to easy scalability of attacks!**

# Plug-and-Secure Communication for CAN Attacker Model with Physical Access to CAN Bus



- Physical access enables more sophisticated attacks (e.g., exploitation of timing or attenuation effects)
- Attacker needs detailed knowledge of the CAN bus

**BUT:**

- With physical access, an attacker could compromise a vehicle much easier (e.g., cut a cable)
- Attacks requiring physical access do not scale; threat with physical access always existed
- Countermeasures are possible → e.g., artificial (random) jitter in bit timing

# Plug-and-Secure Communication for CAN

## Remote Attacks



**Idea: Passively eavesdrop on the channel during key setup**

- ▶ Fact: Alice + Bob derive secret only from secure bit pairs (“00” on bus)



A passive Eve cannot determine the established secret bits



**Idea: Actively interfere with key establishment procedure**

- Action: Eve transmits dominant bits → leads to “00” bit pairs on bus
- Result: Alice + Bob derive different secrets
- Solution: Perform key verification after key generation



An active Eve can prevent a successful key establishment

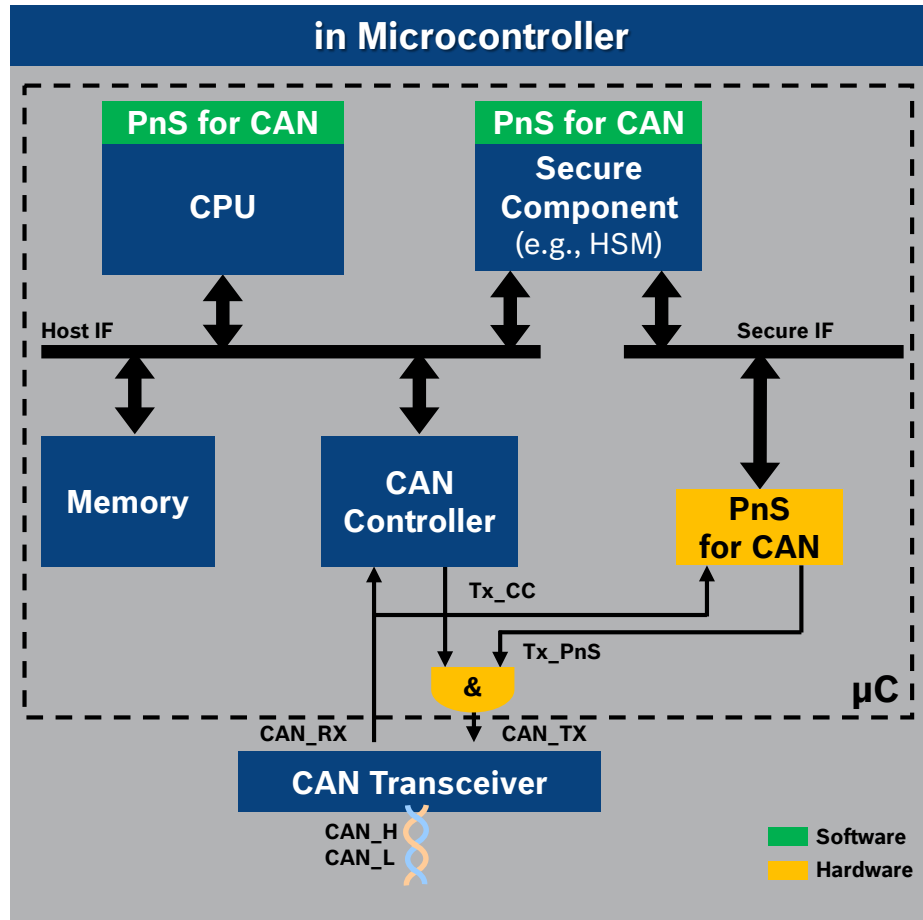


An active Eve cannot determine or influence the established keys



# Plug-and-Secure Communication for CAN

## Implementation: PnS Module in Microcontroller

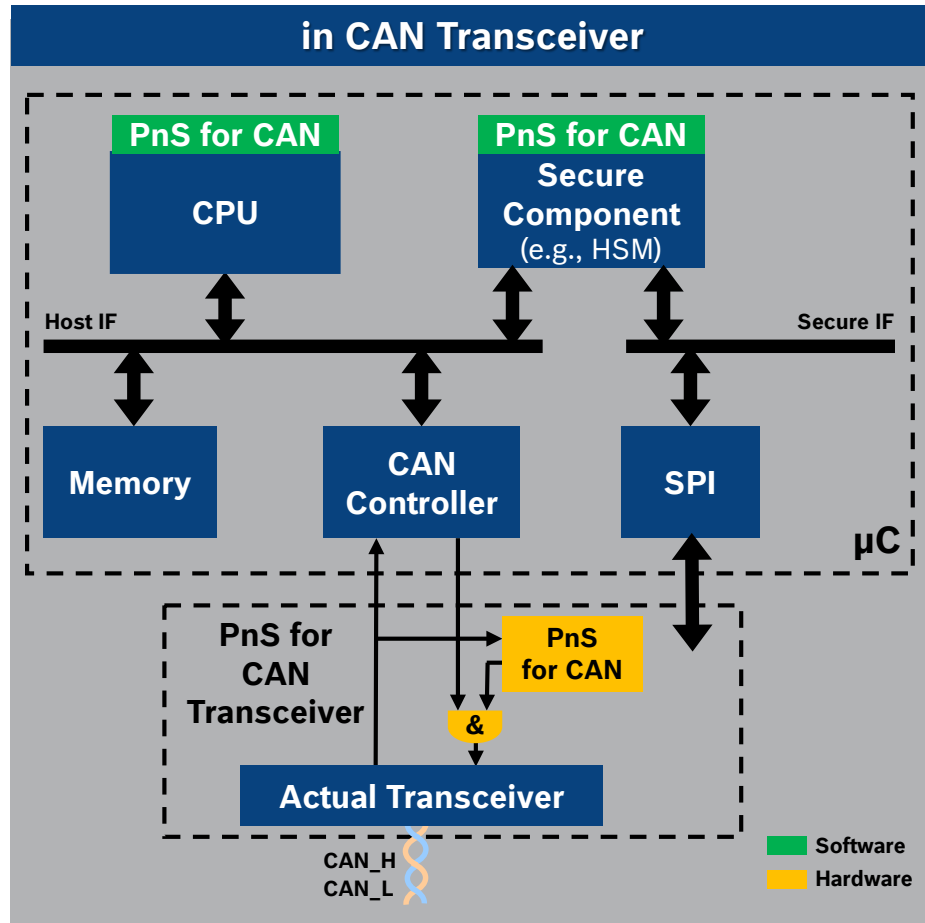


- ### Properties
- ❑ “PnS for CAN” module
    - ❑ is a reduced CAN controller → low costs
    - ❑ connects to the CAN bus in parallel to CAN controller
    - ❑ competes with the on-chip CAN controller and other CAN devices via arbitration
    - ❑ **is compatible to any CAN controller**
  - ❑ Separation of core functions in dedicated HW module is good from a security point of view
  - ❑ Secure component (optional) stores keys and performs crypto functions<sup>1</sup>

<sup>1</sup>If no secure component is available, “PnS for CAN” module might be directly connected to CPU

# Plug-and-Secure Communication for CAN

## Implementation: PnS Module in CAN Transceiver

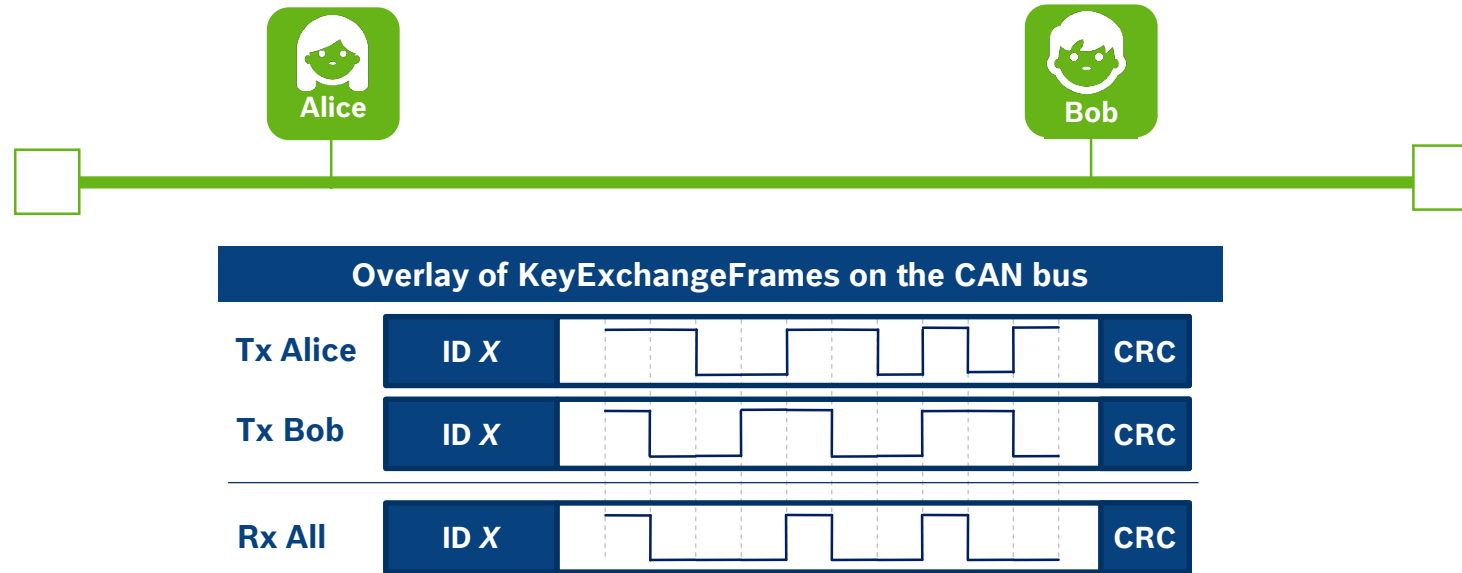


### Properties

- ❑ No modifications of existing µC HW necessary  
→ quick upgrade path
- ❑ May be combined with any existing µC
- ❑ Communication between “PnS for CAN” module in TRX and µC via SPI
- ❑ Encapsulation of core functions in HW module good from a security point of view
- ❑ Secure component (optional) stores keys and performs crypto functions<sup>1</sup>

# Plug-and-Secure Communication for CAN

## Target: KeyExchangeFrames of Alice and Bob need to overlay



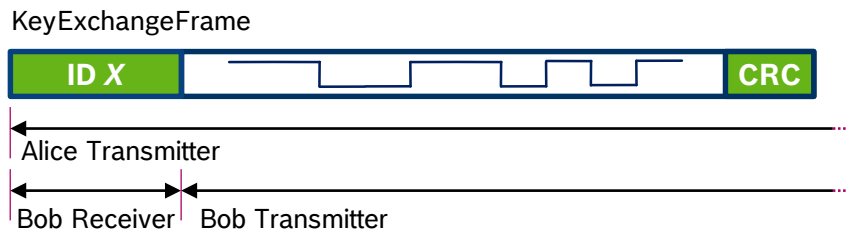
- ▶ **Challenge:** Alice and Bob need to synchronize
- ▶ **Therefore:** One node triggers the other node  
E.g.: Alice tells Bob to start TX of KeyExchangeFrame

# Plug-and-Secure Communication for CAN

## Synchronization of Frame Transmission

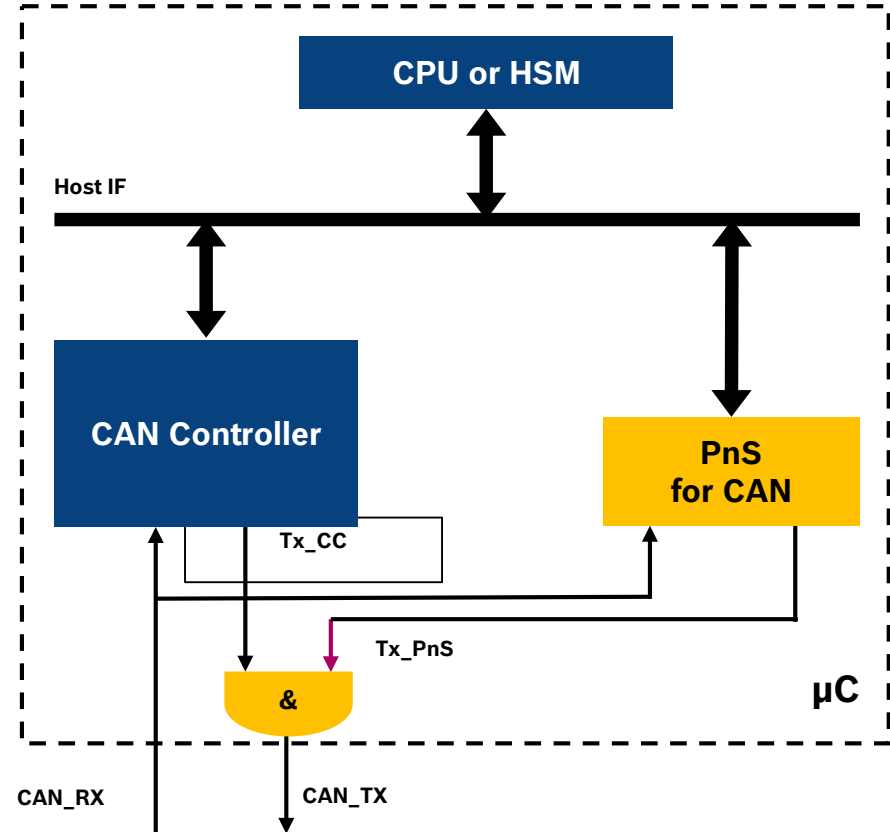
### Procedure

- ▶ Alice/Bob: CPU configures frame ID  $X$  in PnS
- ▶ Alice: Sends a frame with ID  $X$
- ▶ Bob: When PnS detects frame ID  $X$  on CAN bus, it switches its status from Receiver to Transmitter



### Properties

Independent, precise, simple



# Plug-and-Secure Communication for CAN Summary

## ▶ Task

Secure Key Establishment

## ▶ Properties

Very low complexity, high efficiency, low cost



## ▶ Operation

On any CAN bus (Classical CAN or CAN FD)



## ▶ Implementation

PnS Module required in Microcontroller or Transceiver

PnS  
for CAN

## ▶ Major Strengths

Remote / SW-based attacks, Automated key exchange

